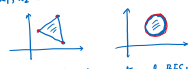


- Last time:
- Thresholding based methods
- General diversity measures/convex optimization.
- Today:
- Extreme points, BFS, concave opt.
- Properties of local optima
- Rewriting based methods.

Recall: $\min_{x \in \mathbb{R}^n} g(x)$ s.t. $Ax = y$
Separable: $g(x) = \sum_{i=1}^n g_i(x_i)$
 $g_i(\cdot)$ monotonically \uparrow in $|x_i|$, $g_i(0) = 0$.
Examples: $g(x) = |x|$, $x \geq 0$
 $g(x) = \ln(x+1)$, $x \geq 0$.

Reformulation:
(P1) $\min_{x \in \mathbb{R}^n} J(x) = \sum_{i=1}^n g_i(x_i)$ s.t. $Ax = y$
(P2) Define $w = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} \in \mathbb{R}^n$
 $\min_{w \in \mathbb{R}^n} J(w) = \sum_{i=1}^n [g_i(x_i) + g_i(-x_i)] - g_i(0)$
s.t. $[A \ -A]w = y$, $w \geq 0$.

(H1): There is a 1-1 mapping betw the local minima of (P1) and (P2), and the value of the objective fun. is the same at corresponding local minima.
Basic feasible soln (BFS): a soln. s.t. the cols of $[A \ -A]$ corresponding to $w_i > 0$ are LI.
Extreme pt.: A point x in a convex set \mathcal{C} is said to be an extreme pt. of \mathcal{C} if there are no distinct $x_1, x_2 \in \mathcal{C}$ s.t. $x = \alpha x_1 + (1-\alpha)x_2$ for some $\alpha \in (0,1)$.



Equivalence of extreme pts and BFS:
Let \mathcal{K} be the convex polytope consisting of all N -vec. s.t. $Ax = y, x \geq 0$.
A vec. $x \in \mathcal{K}$ is an extreme pt. of \mathcal{K} iff x is a BFS to \mathcal{C} .
Proof: Suppose x is a BFS, and x_1, \dots, x_k are nonzero. Then a_1, \dots, a_k are LI. If x is not an extreme pt., then $\exists u, v \in \mathcal{K}$ s.t. $x = \alpha u + (1-\alpha)v$, $0 < \alpha < 1$, $u \neq v$, $u \geq 0, v \geq 0$.

$$\begin{bmatrix} x_1 \\ \vdots \\ x_k \\ \vdots \\ x_n \end{bmatrix} = \alpha \begin{bmatrix} u_1 \\ \vdots \\ u_k \\ \vdots \\ u_n \end{bmatrix} + (1-\alpha) \begin{bmatrix} v_1 \\ \vdots \\ v_k \\ \vdots \\ v_n \end{bmatrix}$$

$\Rightarrow u_{k+1} = \dots = u_n = v_{k+1} = \dots = v_n = 0$.
 $u, v \in \mathcal{K} \Rightarrow Au = y, Av = y$
 $a_1(u_1 - v_1) + \dots + a_k(u_k - v_k) = 0$.
 $u \neq v \Rightarrow a_1, \dots, a_k$ are LD (contradiction).
Conversely, suppose x is an extreme pt. of \mathcal{K} and x_1, \dots, x_k are nonzero. If a_1, \dots, a_k are LD, then $\exists u = \begin{bmatrix} u_1 \\ \vdots \\ u_k \\ 0 \end{bmatrix}$ s.t. $Au = 0$.
Since $x_1, \dots, x_k > 0$, can select $\epsilon > 0$ s.t.
 $\hat{x}_1 + \hat{x}_2, \hat{x}_1 - \hat{x}_2 \in \mathcal{K}$, and
 $x = \frac{1}{2} \hat{x}_1 + \frac{1}{2} \hat{x}_2$
 $\Rightarrow x$ is not an extreme pt. (contradiction).
Hence, x is a BFS.

- Remarks:
1. If the convex set \mathcal{K} corresponds to \mathcal{C} ($Ax=y, x \geq 0$) is nonempty, it has at least one extreme pt.
2. If the LP has a finite optimal soln., it has a finite optimal soln. which is an extreme pt.
3. \mathcal{K} has at most a finite # extreme pts.
4. If the convex polytope \mathcal{K} is full, then \mathcal{K} is a convex polyhedron, i.e., \mathcal{K} consists of points that are a convex comb. of a finite # pts.

Local minima of concave fun.
Let $g(x) = \sum_{i=1}^n g_i(x_i)$, strictly concave for $x \geq 0$.
(P3) $\min_{x \in \mathbb{R}^n} g(x)$ s.t. $Ax = y, x \geq 0$.
Claim: When $g(\cdot)$ is strictly concave, all local minima of (P3) occur at extreme pts of \mathcal{K} .
 \Rightarrow all local minima occur at BFS to \mathcal{C} .
Proof: Let x_0 be a local min.
 $Ax_0 = y, x_0 \geq 0, \exists \epsilon > 0$ s.t. $g(x_0) \leq g(x')$ for $x' \in \mathcal{C}, \|x_0 - x'\|_2 \leq \epsilon$.
If x_0 is not an extreme pt., then $\exists x_1, x_2 \in \mathcal{K}, x_1 \neq x_2, \lambda \in (0,1)$ s.t. $x_0 = \lambda x_1 + (1-\lambda)x_2$.
By the strict concavity of $g(\cdot)$,
 $g(x_0) > \lambda g(x_1) + (1-\lambda)g(x_2) \geq \min(g(x_1), g(x_2))$.
Now, pick any $0 \neq \delta \in \mathcal{N}(A), \text{supp}(\delta) \subseteq \text{supp}(x_0)$.
Then, for small enough δ , $x_0 \pm \delta$ is feasible.
 $x_0 = \frac{1}{2}(x_0 + \delta) + \frac{1}{2}(x_0 - \delta)$
 $\Rightarrow g(x_0) > \frac{1}{2}g(x_0 + \delta) + \frac{1}{2}g(x_0 - \delta) \geq \min(g(x_0 + \delta), g(x_0 - \delta))$
 $\Rightarrow g(x_0) > \min(g(x_0 + \delta), g(x_0 - \delta))$
 $\Rightarrow (x_0 \pm \delta)$ is a better pt. than x_0 (contradiction). \square

General result on maximizing concave fun:

Then, let f be a convex fn. defined on the bounded, closed convex set S . If f has a maximum over S , it is achieved at an extreme pt. of S .

Cor. The global min. of (P_1) or (P_2) is also a PFS.

$$\min_{x \in \mathbb{R}^n} \sum_{i=1}^m g_i(x) \quad \text{s.t.} \quad Ax=y, \quad x \geq 0.$$

1. $g(x) = g(-x) = g(x)$
2. $g(x)$ is monotone \uparrow for $x \in \mathbb{R}_+$
3. $g(x)$ is strictly concave $x \in \mathbb{R}_+$.

Naive data:

$$\min_{\substack{x \in \mathbb{R}^n \\ x \geq 0 \\ x \geq 0}} \sum_{i=1}^m g_i(x_i) + \|Ax-y\|_2$$

$$\min_{\substack{x \in \mathbb{R}^n \\ x \geq 0 \\ x \geq 0}} \sum_{i=1}^m g_i(x_i) \quad \text{s.t.} \quad \|Ax-y\|_2 \leq \epsilon.$$

A general approach: Majorization-minimization (MM):

Let $f(x)$ be a fn. to be minimized.
 Let $\underline{g}(x|o^{(k)}) \geq f(x) \neq 0$ $\left. \begin{array}{l} \underline{g}(o^{(k)}) \text{ as} \\ \text{upper bd. on} \\ f(x), \forall x \\ \text{at } o^{(k)}. \end{array} \right\}$
 ($\underline{g}(o^{(k)})$) $\underline{g}(o^{(k)}|o^{(k)}) = f(o^{(k)})$.

MM algo:
 Init. $o^{(1)}$ = something convenient.

Iterate $o^{(k+1)} = \arg \min_x \underline{g}(x|o^{(k)})$ — Compute the arg. min. numerically.
 determine $\underline{g}(o^{(k+1)})$.

Until convergence.
 $f(o^{(k+1)}) \leq \underline{g}(o^{(k+1)}|o^{(k)}) \leq \underline{g}(o^{(k)}|o^{(k)}) = f(o^{(k)})$